

# Lanczos-like method for the time-ordered exponential

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# Time-ordered exponential

Let  $t \in I \subseteq \mathbb{R}$ ,  $A(t)$  a time dependent matrix. The **time-ordered exponential** is the unique solution  $U(t)$  of the system of coupled linear differential equations

$$A(t)U(t) = \frac{d}{dt}U(t), \quad U(0) = I.$$

If  $A(t')A(t) = A(t)A(t')$  for all  $t, t'$ , then

$$U(t) = \exp\left(\int_0^t A(\tau) d\tau\right).$$

Otherwise,  $U$  has **no known explicit form in terms of  $A$** . It can be denoted by the time-ordering operator  $\mathcal{T}$  as

$$U(t) = \mathcal{T} \exp\left(\int_0^t A(\tau) d\tau\right).$$

[Dyson, 1952]

# Time-ordered exponential

The time-ordering expression is more a notation than an explicit form as the action of the time-ordering operator is very difficult to evaluate.

- **Applications:** system dynamics (quantum physics); see, e.g., [Blanes & al., 2009].
- **Classical approaches** Magnus series [Magnus, 1954]: infinite series of increasing complexity (almost always divergent).  
Using Floquet theory: only  $A$  periodic, the solution is given in terms of a infinite system. See, e.g., [Blanes & al., 2009].
- **Path-sum approach:** based on the *path-sum* decomposition of the graph whose  $A$  is the adjacency matrix [Giscard & al., 2015]. The expression has a finite number of terms, but it may be too complicate.

- Idea: model reduction + path-sum approach;
- Lanczos algorithm properties (time-independent matrices);
- \*-Lanczos algorithm for time-ordered exponentials.

Warning 1: Work in progress!

Comments and suggestions are more than welcome!

Warning 2: No computational approach

We will derive an expression for the solution on the base of which a further numerical approach may be developed.

Let  $t' \geq t$  and let  $A_1(t', t)$  and  $A_2(t', t)$  be (doubly) time-dependent matrices. We define the **convolution-like product**

$$(A_2 * A_1)(t', t) := \int_t^{t'} A_2(t', \tau) A_1(\tau, t) d\tau.$$

Consider the **\*-resolvent**

$$R_*(A) := (1_* - A)^{* - 1} = 11_* + \sum_{k \geq 1} A^{*k}$$

(it exists if every entry of  $A$  is bounded for all  $t', t$ ). Then

$$U(t', t) = \int_t^{t'} R_*(A)(\tau, t) d\tau,$$

with  $1_* = \delta(t' - t)$  the Dirac delta function; [Giscard & al., 2015].

# Path-sum formulation

If we look at  $A$  as the adjacency matrix of the graph  $G$  with time-dependent weights, then the  $(i, j)$  element of the  $*$ -resolvent,  $(R_*)_{ij}$ , can be rewritten in terms of the **simple paths** from  $i$  to  $j$  in  $G$ ; [Giscard & al., 2015].

This approach provides **exact solutions** and it is unconditionally convergent. However, it suffers from drawbacks:

- It requires to find all the simple cycles and simple paths of the graph  $G$  (computationally expensive);
- The expression for  $U$  may be too complicated due to the structure of the graph  $G$ .

We want to compute a matrix  $T$  so that

- $T$  is **tridiagonal**, i.e., it correspond to a graph which is a path (with self-loops);
- $R_*(T)_{11} = R_*(A)_{jj}$ , i.e., they have the same **\*-moments**

$$(T^{*k})_{11} = (A^{*k})_{jj}, \quad j = 0, 1, 2, \dots$$

In the simple case in which  $A$  is time-**independent**, a  $T$  with analogous properties is give as the output of the last iteration of **Lanczos algorithm**.

We will derive a Lanczos-like method based on an **\*-orthogonalization process**.

# Non-Hermitian Lanczos algorithm

Let  $A$  be a time-independent matrix and  $\mathbf{v}, \mathbf{w}$  time-independent vectors. Our goal is to approximate

$$\mathbf{w}^H \exp(A) \mathbf{v}.$$

Consider the **Krylov subspaces**

$$\text{span}\{\mathbf{v}, A\mathbf{v}, \dots, A^{n-1}\mathbf{v}\}, \quad \text{span}\{\mathbf{w}, A^H\mathbf{w}, \dots, (A^H)^{n-1}\mathbf{w}\}.$$

Assuming no breakdowns, the **non-Hermitian Lanczos** computes the matrices

$$V_n = [\mathbf{v}_0, \dots, \mathbf{v}_{n-1}], \quad W_n = [\mathbf{w}_0, \dots, \mathbf{w}_{n-1}]$$

respectively basis of the Krylov subspaces, so that  $W_n^H V_n = I$ .



# Non-Hermitian Lanczos algorithm

The tridiagonal matrix defined as

$$T_n = W_n^H A V_n,$$

satisfies the **matching moment property**

$$\mathbf{w}^H A^k \mathbf{v} = \mathbf{e}_1^H T_n^k \mathbf{e}_1, \quad k = 0, \dots, 2n - 1.$$

Hence we get the approximation (**model reduction**)

$$\mathbf{w}^H \exp(A) \mathbf{v} \approx \mathbf{e}_1^H \exp(T_n) \mathbf{e}_1;$$

e.g., [Golub, Meurant, 2010].

In particular, for  $n$  equal to the size of  $A$  (or for lucky or incurable breakdowns) we get **exactness**.

## \*-resolvent approximation

Now,  $A(t)$  is **time-dependent** and we want to approximate the  $i, j$  element of the time-ordered exponential of  $A$ , i.e.,

$$U(t', t)_{i,j} = \mathbf{e}_i^H \mathcal{T} \exp \left( \int_t^{t'} A(\tau) d\tau \right) \mathbf{e}_j = \int_t^{t'} \mathbf{e}_i^H R_*(A)(\tau, t) \mathbf{e}_j d\tau.$$

Hence we want to approximate the value

$$\mathbf{w}^H R_*(A)(t', t) \mathbf{v},$$

with  $\mathbf{v}, \mathbf{w}$  **time-independent** vectors.

# Time-dependent Krylov subspaces

Given the distributions  $\alpha_0(t', t), \dots, \alpha_k(t', t)$  (**coefficients**),

**matrix \*-polynomial:** 
$$p(A)(t', t) := \sum_{j=0}^k (A^{*j} * \alpha_j)(t', t),$$

**dual \*-polynomial:** 
$$p^D(A)(t', t) := \sum_{j=0}^k (\bar{\alpha}_j * (A^{*j})^H)(t', t),$$

and hence the **time-dependent Krylov subspaces** as

$$\mathcal{K}_n(A, \mathbf{v})(t', t) := \{ (p(A)\mathbf{v}) \mid p \text{ of degree } \leq n-1 \},$$
$$\mathcal{K}_n^D(A^H, \mathbf{w})(t', t) := \left\{ \left( \mathbf{w}^H p^D(A^H) \right) \mid p \text{ of degree } \leq n-1 \right\}.$$

# Building up the \*-Lanczos process

We aim to build basis for  $\mathcal{K}_n(A, \mathbf{v})$  and  $\mathcal{K}_n^D(A^H, \mathbf{w})$

$$V_n = [\mathbf{v}_0, \dots, \mathbf{v}_{n-1}], \quad W_n^H = [\mathbf{w}_0, \dots, \mathbf{w}_{n-1}]^H$$

which are \*-biorthonormal, i.e.,

$$W_n^H * V_n = I_{1*}.$$

Assuming  $\mathbf{w}^H \mathbf{v} = 1$ , the first vectors are

$$\mathbf{v}_0 = \mathbf{v} 1_*, \quad \mathbf{w}_0^H = \mathbf{w}^H 1_*.$$

# Building up the \*-Lanczos process

Consider a vector  $\widehat{\mathbf{v}}_1 \in \mathcal{K}_2(A, \mathbf{v})$

$$\widehat{\mathbf{v}}_1 = A * \mathbf{v}_0 - \mathbf{v}_0 * \alpha_0 = A\mathbf{v} - \mathbf{v}\alpha_0.$$

Under the condition  $\mathbf{w}_0^H * \widehat{\mathbf{v}}_1 = 0$ , we get

$$\alpha_0 = \mathbf{w}_0^H * A * \mathbf{v}_0 = \mathbf{w}^H A \mathbf{v} \quad (1)$$

Similarly,

$$\widehat{\mathbf{w}}_1^H = \mathbf{w}_0^H * A - \alpha_0 * \mathbf{w}_0^H = \mathbf{w}^H A - \alpha_0 \mathbf{w}^H,$$

with  $\alpha_0$  given by (1).

# Building up the \*-Lanczos process

The condition

$$\mathbf{w}_1^H * \mathbf{v}_1 = 1_*,$$

is satisfied defining

$$\mathbf{v}_1 = \widehat{\mathbf{v}}_1 * (\beta_1)^{*-1}, \quad \mathbf{w}_1 = \widehat{\mathbf{w}}_1,$$

with

$$\beta_1 = \widehat{\mathbf{w}}_1^H * \widehat{\mathbf{v}}_1 = \mathbf{w}^H \mathbf{A}^{*2} \mathbf{v} - \alpha_0^{*2}.$$

We assume  $\beta_1$  \*-invertible and computable. We will discuss this issue later.

# Building up the \*-Lanczos process

By induction, consider the vector

$$\widehat{\mathbf{v}}_n = \mathbf{A} * \mathbf{v}_{n-1} - \sum_{i=0}^{n-1} \mathbf{v}_i * \gamma_i.$$

The condition  $\mathbf{w}_j^H * \widehat{\mathbf{v}}_n = 0$  for  $j = 0, \dots, n-1$  gives

$$\gamma_j = \mathbf{w}_j^H * \mathbf{A} * \mathbf{v}_{n-1}, \quad j = 0, \dots, n-1.$$

Since  $\mathbf{w}_j^H * \mathbf{A} \in \mathcal{K}_{j+1}^D(\mathbf{A}^H, \mathbf{w})$  we get

$$\gamma_j = 0, \quad j = 0, \dots, n-3.$$

# Building up the \*-Lanczos process

Assuming  $\mathbf{v}_{-1} = \mathbf{w}_{-1} = 0$ , we get the three-term recurrences

$$\begin{aligned}\mathbf{w}_n^H &= \mathbf{w}_{n-1}^H * \mathbf{A} - \alpha_{n-1} * \mathbf{w}_{n-1}^H - \beta_{n-1} * \mathbf{w}_{n-2}^H, \\ \mathbf{v}_n * \beta_n &= \mathbf{A} * \mathbf{v}_{n-1} - \mathbf{v}_{n-1} * \alpha_{n-1} - \mathbf{v}_{n-2},\end{aligned}$$

with the coefficients given by

$$\alpha_{n-1} = \mathbf{w}_{n-1}^H * \mathbf{A} * \mathbf{v}_{n-1}, \quad \beta_n = \mathbf{w}_n^H * \mathbf{A} * \mathbf{v}_{n-1}.$$

- If  $\beta_n$  is not \*-invertible we have a **breakdown**;
- We assume no breakdowns in this first study.



# Building up the \*-Lanczos process

For  $n \geq 2$ , setting  $\gamma_n = (\mathbf{w}_n^H + \mathbf{w}_{n-2}^H) * A * \mathbf{v}_{n-1}$

$$\beta_n = \gamma_n - 1_*, \quad \beta_n^{*-1} = R_*(\gamma_n),$$

which can be reformulated as a **Volterra equations of the second kind**

$$R_* = 1_* + \gamma_n * R_*.$$

If  $\gamma_n$  is bounded over  $I^2$ , then  $R_*$  necessarily exists and the Neumann series  $\sum_{k \geq 0} \gamma_n^{*k}$  converges to it super-geometrically.

## \*-inverse, one time

Let  $a(t') \neq 0$  be a function of a single time variable  $t' \in I$ . Then  $a^{*-1}$  exists as a distribution and its given by

$$a^{*-1}(t', t) = \frac{d}{dt'} \left[ \frac{\delta(t' - t)}{a(t')} \right] = \frac{\delta'(t' - t)}{a(t')} - \frac{a'(t')}{a(t')^2} \delta(t' - t).$$

$\delta'$  and  $a'$  are the derivatives of the Dirac delta distribution and  $a$ , respectively. Now let  $f(t', t)$  be a distribution. Then

$$\begin{aligned}(a^{*-1} * f)(t', t) &= \frac{d}{dt'} \left[ \frac{f(t', t)}{a(t')} \right] \\(f * a^{*-1})(t', t) &= -\frac{1}{a(t)} \frac{d}{dt} [f(t', t)].\end{aligned}$$

# \*-Lanczos Algorithm

Initialize:  $\mathbf{v}_{-1} = \mathbf{w}_{-1} = \mathbf{0}$ ,  $\mathbf{v}_0 = \mathbf{v} \mathbf{1}_*$ ,  $\mathbf{w}_0^H = \mathbf{w}^H \mathbf{1}_*$ .

$$\alpha_0 = \mathbf{w}^H \mathbf{A} \mathbf{v},$$

$$\mathbf{w}_1^H = \mathbf{w}^H \mathbf{A} - \alpha_0 \mathbf{w}^H,$$

$$\widehat{\mathbf{v}}_1 = \mathbf{A} \mathbf{v} - \mathbf{v} \alpha_0,$$

$$\beta_1 = \mathbf{w}^H \mathbf{A}^2 \mathbf{v} - \alpha_0^2,$$

Determine  $\beta_1^{*-1}$ ,

$$\mathbf{v}_1 = \widehat{\mathbf{v}}_1 * \beta_1^{*-1},$$

For  $n = 2, \dots$

$$\alpha_{n-1} = \mathbf{w}_{n-1}^H * \mathbf{A} * \mathbf{v}_{n-1},$$

$$\mathbf{w}_n^H = \mathbf{w}_{n-1}^H * \mathbf{A} - \alpha_{n-1} * \mathbf{w}_{n-1}^H - \beta_{n-1} * \mathbf{w}_{n-2}^H,$$

$$\widehat{\mathbf{v}}_n = \mathbf{A} * \mathbf{v}_{n-1} - \mathbf{v}_{n-1} * \alpha_{n-1} - \mathbf{v}_{n-2},$$

$$\beta_n = \mathbf{w}_n^H * \mathbf{A} * \mathbf{v}_{n-1},$$

If  $\beta_n$  is not \*-invertible, then stop, otherwise,

$$\mathbf{v}_n = \widehat{\mathbf{v}}_n * \beta_n^{*-1},$$

end.

# Tridiagonal matrix

$$T_n := \begin{bmatrix} \alpha_0 & 1_* & & & \\ \beta_1 & \alpha_1 & \ddots & & \\ & \ddots & \ddots & & \\ & & & \ddots & 1_* \\ & & & \beta_{n-1} & \alpha_{n-1} \end{bmatrix},$$

and recall  $V_n := [\mathbf{v}_0, \dots, \mathbf{v}_{n-1}]$  and  $W_n := [\mathbf{w}_0, \dots, \mathbf{w}_{n-1}]$ . Then

$$\begin{aligned} A * V_n &= V_n * T_n + (\mathbf{v}_n * \beta_n) \mathbf{e}_n^H \\ W_n^H * A &= T_n * W_n^H + \mathbf{e}_n \mathbf{w}_n^H. \end{aligned}$$

Hence  $T_n$  is the  $*$ -projection of  $A$  onto  $\mathcal{K}_n(A, \mathbf{v})$  along the direction of  $\mathcal{K}_n^D(A^H, \mathbf{w})$ , i.e.,

$$T_n = W_n^H * A * V_n.$$

# Matching moment

## Matching moment property

Let  $A$ ,  $\mathbf{w}$ ,  $\mathbf{v}$  and  $T_n$  be as described above, then

$$\mathbf{w}^H(A^{*k})\mathbf{v} = \mathbf{e}_1^H(T_n^{*k})\mathbf{e}_1, \quad \text{for } k = 0, \dots, 2n - 1.$$

In particular, for  $n$  equal to the size of  $A$  we get

$$\mathbf{w}^H(A^{*k})\mathbf{v} = \mathbf{e}_1^H(T_n^{*k})\mathbf{e}_1, \quad \text{for } k = 0, 2, \dots,$$

and, hence,

$$\mathbf{w}^H R_*(A)\mathbf{v} = \mathbf{e}_1^H R_*(T_n)\mathbf{e}_1.$$

# Idea of the proof

Inspired by the classical Lanczos algorithm, we follow the analogy with **formal orthogonal polynomial** properties; see [Draux, 1982], also the surveys [Pozza & al., 2017, 2018].

Define the set of  $*$ -polynomials

$$\mathcal{P}_* := \left\{ p(\lambda) = \sum_{j=0}^k \lambda^{*j} * \gamma_j(t', t) \right\},$$

with coefficients the  $\gamma_j(t', t)$  distributions.

Consider a  $*$ -sesquilinear form  $[\cdot, \cdot] : \mathcal{P}_* \times \mathcal{P}_* \rightarrow \mathbb{C}[I^2]$

$$[q * \alpha, p * \beta] = \bar{\alpha} * [q, p] * \beta,$$

$$[q_1 + q_2, p_1 + p_2] = [q_1, p_1] + [q_2, p_1] + [q_1, p_2] + [q_2, p_2].$$

# Idea of the proof

The polynomials  $p_0, p_1, \dots$  and  $q_0, q_1, \dots$  are  $*$ -biorthonormal with respect to  $[\cdot, \cdot]$ , if  $[q_i, p_j] = \delta_{ij} 1_*$ .

## Lemma

Let  $p_0, \dots, p_{n-1}$  and  $q_0, \dots, q_{n-1}$  be  $*$ -biorthonormal with respect to  $[\cdot, \cdot]_A$  and to  $[\cdot, \cdot]_B$ . If  $[1_*, 1_*]_A = [1_*, 1_*]_B = 1_*$ , then

$$[\lambda^{*k}, 1_*]_A = [\lambda^{*k}, 1_*]_B, \quad k = 0, \dots, 2n - 1.$$

Finally we showed that the forms

$$[q, p] := \mathbf{w}^H \bar{q}^D (A^H) * p(A) \mathbf{v}_., \quad [q, p]_n := \mathbf{e}_1^H \bar{q}^D (T_n^H) * p(T_n) \mathbf{e}_1,$$

have the same first  $n$   $*$ -biorthonormal polynomials.

## \*-Lanczos approximation

Consider the time-ordered exponential  $U_n$  given by

$$T_n(t)U_n(t) = \frac{d}{dt}U_n(t).$$

The **matching moment property** justify the use of the approximation

$$\mathbf{w}^H U(t) \mathbf{v} \approx \mathbf{e}_1^H U_n(t) \mathbf{e}_1 = \mathbf{e}_1^H \int_0^t R_*(T_n)(\tau, 0) d\tau \mathbf{e}_1.$$

The **path-sum method** gives the finite **\*-continued fraction**

$$\left( R_*(T_n)(\tau, 0) \right)_{11} = \left( 1_* - \alpha_0 - (1_* - \alpha_1 - (\dots_* \beta_{n-1})^{*-1} \dots_* \beta_2)^{*-1} \dots_* \beta_1 \right)^{*-1},$$

see [Giscard & al. 2012, 2015].



## \*-Lanczos approximation

- The method converges at most in  $N$  iteration,  $N$  being the size of  $A$ ;
- The intermediate approximations  $U_n$  have a clear meaning in terms of moment matching;
- An explicit expression for  $\beta_n^{*-1}$  may not be easily computed (work in progress).

# Error bound?

We have not yet derived an a priori bound for the error

$$E_n(t) = \int_0^t \sum_{k=2n}^{\infty} \mathbf{w}^H A^{*k}(\tau) \mathbf{v} - \mathbf{e}_1^H T_n^{*k}(\tau, 0) \mathbf{e}_1 d\tau.$$

However, notice that if it is possible to bound

$$\sup_{t \in I} |\mathbf{w}^H A(t) \mathbf{v}| \leq C, \quad \sup_{t \in I} |\mathbf{e}_1^H T_n(t, 0) \mathbf{e}_1| \leq D,$$

then we get

$$|E_n(t)| \leq \frac{((C + D)t)^{2n+1} e^{(C+D)t}}{(2n + 1)!}.$$

Note that

$$\left| \int_0^t A^{*2n}(\tau) d\tau \right| \leq C^{*2n} \int_0^t d\tau = \frac{(Ct)^{2n+1}}{(2n + 1)!}.$$

# Breakdown

When  $\beta_n$  is not invertible we have a **breakdown**. We distinguish two cases

- **Lucky breakdown**: if either  $\mathbf{w}_n \equiv 0$  or  $\mathbf{v}_n \equiv 0$ ;
- **Serious breakdown**: if  $\mathbf{v}_n, \mathbf{w}_n \neq 0$ , and  $\mathbf{w}_n^H * \mathbf{v}_n$  is not invertible.

## Proposition

If  $\mathbf{v}_n = 0$  (or  $\mathbf{w}_n = 0$ ), then

$$\mathbf{w}^H(A^{*k})\mathbf{v} = \mathbf{e}_1^H(T_n^{*k})\mathbf{e}_1, \quad \text{for } k = 0, 1, \dots$$

Hence

$$\mathbf{w}^H U(t)\mathbf{v} = \mathbf{e}_1^H \int_0^t \left(1_* - T_n\right)^{* - 1}(\tau, 0) d\tau \mathbf{e}_1.$$

# Serious breakdown

- We do not treat the possible solutions to a general serious breakdowns;
- A possible way to deal with it is to use a **look-ahead strategy** analogous to what can be done for the non-Hermitian Lanczos algorithm.

## A particular and relevant case of serious breakdown.

If  $A$  is a sparse non-Hermitian matrix, then it may be that  $A_{ij} = 0$  and  $A^{*2}_{ij} = 0$ . Then  $*$ -Lanczos with inputs  $A, \mathbf{e}_i, \mathbf{e}_j$  gives  $\beta_1 = 0$ . To fix this issue, we can consider the problem

$$\mathbf{e}_i^H U(t) \mathbf{e}_j = (\mathbf{e} + \mathbf{e}_i)^H U(t) \mathbf{e}_j - \mathbf{e}^H U(t) \mathbf{e}_j,$$

with  $\mathbf{e} = (1, \dots, 1)^H$ . Then one can approximate the  $(\mathbf{e} + \mathbf{e}_i)^H U \mathbf{e}_j$  and  $\mathbf{e}^H U \mathbf{e}_j$  separately, which now are less likely going to have a breakdown at the first step since  $\mathbf{e}$  is full; e.g., [Golub, Meurant, 2010].

## Example: Ordinary matrix exponential

Consider a constant matrix

$$A = \begin{pmatrix} -1 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & -1 \end{pmatrix}.$$

Since  $A$  commutes with itself at all times,  $\mathcal{T}e^{\int A(\tau)d\tau} \equiv e^{A(t)}$ .

$$(e^{A(t)})_{11} = -\frac{1}{2} \sinh(2t) + \frac{1}{2} \cosh(2t) + \frac{1}{2} \cosh(\sqrt{2}t).$$

Note that we chose a symmetric matrix for simplicity.

## Example: Ordinary matrix exponential

We compute the result by \*-Lanczos. We set

$$\mathbf{w}^H := \mathbf{v}^H := (1, 0, 0), \quad \mathbf{w}_0 = \mathbf{w}1_*, \quad \mathbf{v}_0 = \mathbf{v}1_*$$

Note that in such a simple setting we have

$$A^{*n} = A^n \times 1^{*n} = A \times (t' - t)^{n-1} / (n-1)!$$

It follows

$$\alpha_0 = \mathbf{w}^H A \mathbf{v} = -1,$$

$$\mathbf{w}_1^H = \mathbf{w}^H A - \alpha_0 \mathbf{w}^H = \hat{\mathbf{v}}_1^H = (0, 1, 1),$$

$$\beta_1 = \mathbf{w}^H A^2 \mathbf{v} \times (t' - t) - \alpha_0^2 (t' - t) = 2(t' - t),$$

$$\beta_1^{*-1} = \frac{1}{2} \delta''(t' - t).$$

## Example: Ordinary matrix exponential

We rescaled  $\widehat{\mathbf{v}}_1$

$$\mathbf{v}_1 = \widehat{\mathbf{v}}_1 * \beta_1^{*-1} = \frac{1}{2} \delta'(t' - t) \times (0, 1, 1)^H,$$

from which we obtain

$$\alpha_1(t', t) = \mathbf{w}_1 * \mathbf{A} * \mathbf{v}_1 = \frac{1}{2} - \frac{1}{2} \Theta(t - t') = \frac{1}{2},$$

$$\mathbf{w}_2(t', t) = \mathbf{w}_1 * \mathbf{A} - \alpha_1 * \mathbf{w}_1 - \beta_1 * \mathbf{w}_0 = \frac{1}{2} (t' - t) \times (0, 1, -1),$$

$$\widehat{\mathbf{v}}_2(t', t) = \mathbf{A} * \mathbf{v}_1 - \mathbf{v}_1 * \alpha_1 - \mathbf{v}_0 = \frac{1}{4} \delta(t' - t) \times (0, 1, -1)^H,$$

$$\beta_2 = \mathbf{w}_2 * \mathbf{A} * \mathbf{v}_1 = \frac{1}{4} (t' - t).$$

$\Theta(\cdot)$  is the Heaviside theta function ( $\Theta(0) = 1$ ), recall that  $t' \geq t$ .

## Example: Ordinary matrix exponential

$\beta_2$  is \*-invertible with

$$\beta_2^{*-1} = 4\delta''(t' - t),$$

and we get

$$\mathbf{v}_2 = \widehat{\mathbf{v}}_2 * \beta_2^{*-1} = \delta''(t' - t) \times (0, 1, -1)^H,$$

$$\alpha_2 = \mathbf{w}_2 * \mathbf{A} * \mathbf{v}_2 = -\frac{3}{2}.$$



## Example: Ordinary matrix exponential

We have determined the \*-Lanczos matrices  $T$ ,  $V$  and  $W$ .

$$T_3 = \begin{pmatrix} -1 & \delta(t' - t) & 0 \\ 2(t' - t) & \frac{1}{2} & \delta(t' - t) \\ 0 & \frac{1}{4}(t' - t) & -\frac{3}{2} \end{pmatrix},$$

$$V_3 = \begin{pmatrix} \delta(t' - t) & 0 & 0 \\ 0 & \frac{1}{2}\delta'(t' - t) & \delta''(t' - t) \\ 0 & \frac{1}{2}\delta'(t' - t) & -\delta''(t' - t) \end{pmatrix},$$

$$W_3 = \begin{pmatrix} \delta(t' - t) & 0 & 0 \\ 0 & 1 & 1 \\ 0 & \frac{1}{2}(t' - t) & \frac{1}{2}(t' - t) \end{pmatrix}.$$

## Example: Ordinary matrix exponential

Using the continued fraction formulation we get

$$\left( 1_* + 1 - \left( 1_* - \frac{1}{2} - \left( 1_* + \frac{3}{2} \right)^{* - 1} * \frac{1}{4} (t' - t) \right)^{* - 1} * 2(t' - t) \right)^{* - 1}.$$

Since all the entries depend only on  $t' - t$ , we may also use a Laplace transform approach.

$$(\text{Id}_* - \mathbb{T})_{11}^{* - 1}(t, 0) = \sinh(2t) + \frac{1}{\sqrt{2}} \sinh(\sqrt{2}t) - \cosh(2t),$$

which is indeed the derivative of

$$(e^{A(t)})_{11} = -\frac{1}{2} \sinh(2t) + \frac{1}{2} \cosh(2t) + \frac{1}{2} \cosh(\sqrt{2}t).$$

# Conclusion

\*-Lanczos is able to express the solution of a coupled linear differential equations with non-constant coefficient. It presents several difficulties (\*-inverses, integral solutions, ...). Nevertheless,

- It is convergent in a finite number of steps;
- It express the solution with a finite formulation (as the path-sum method);
- It admits an expression in terms of a **simple** finite continued fraction.

Such a general approach has been so far out of reach.

# Future developments and open problems

- Convergence properties of the approximation;
- Explicit formulation of the  $*$ -inverses  $\beta_n^{*-1}$ ;
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Thank you for your attention!