

Weak inverses

for

Strict  $n$ -categories

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\* based on a joint work  
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Notations :

$n$ -Cat

$\mathcal{C}$ ,  $\mathcal{D}$  ..... : (small and) strict  $n$ -categories

$F$ :  $\mathcal{C} \rightarrow \mathcal{D}$  ... : strict  $n$ -functors  
= morphisms

$\alpha$ :  $F \Rightarrow G : \mathcal{C} \rightarrow \mathcal{D}$  ... : weak  $n$ -transform.  
= 2-morphisms

Given  $\mathcal{C}$ , cells :

0-cells : (objects)  $c_0$ ,  $c_0'$ ,  $c_0''$  ...

1-cells : (arrows)  $c_0 \xrightarrow{c_1} c_0$   $c_1'$ , ...

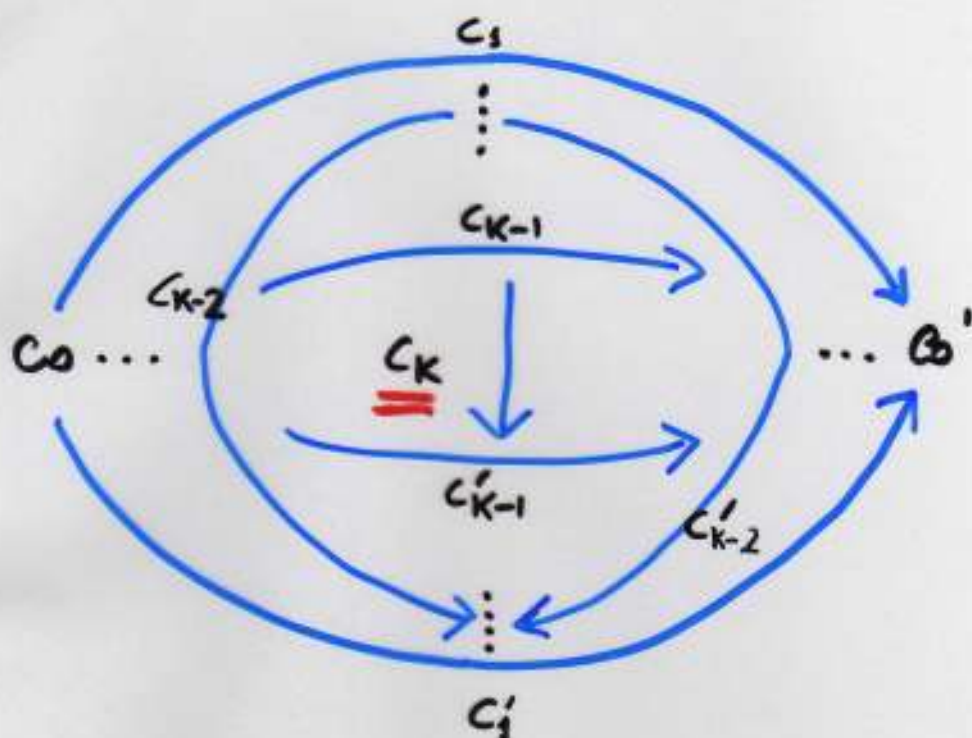
⋮

$k$ -cells :  $c_k : c_{k-1} \rightarrow c_{k-1}'$

( $c_{k-1}, c_{k-1}' : c_{k-2} \rightarrow c_{k-2}' : \dots c_1 \rightarrow c_1' : c_0 \rightarrow c_0'$ )

comp.

## Question



What does it mean for  $C_k$  being invertible?

## Answer

There exists  $C_k^{-1}$  s.t.

$$C_k \circ C_k^{-1} = 1_{C_{k-1}}$$

$$C_k^{-1} \circ C_k = 1_{C'_k}$$

... end of the talk !

If we mean weakly invertible <sup>3</sup>  
then there are several possible answers

I) [R. Street, 1985] (in a  $\omega$ -category)  
defines what means for two  
objects  $c_0, c_0'$  being  $n$ -equivalent:

$\exists$   $c_1: c_0 \rightarrow c_0'$  ,  $c_1^*: c_0' \rightarrow c_0$  s.t.

(i)  $c_1 \circ c_1^*$  and  $1_{c_0}$   $(n-1)$ -equiv. in  $\mathcal{C}(c_0, c_0)$

(ii)  $c_1^* \circ c_1$  and  $1_{c_0'}$   $(n-1)$ -equiv. in  $\mathcal{C}(c_0', c_0')$

of course 0-equiv. = equal.

We will say:

•) equivalent in a  $n$ -cat =  $n$ -equiv.

•)  $c_1$  is  $S$ -invertible if it  
establishes an equivalence  $c_0 \rightsquigarrow c_0'$

## → S-inverses in low dimensions

$\boxed{M=1}$  (Categories)

$c_0 \xrightarrow{c_1} c_0'$  s-invertible if

$$\exists c_1^*: c_0' \rightarrow c_0 \text{ s.t. } 1_{c_0} = c_1 c_1^*$$

$$c_1^* c_1 = 1_{c_0'}$$

→ iso

$\boxed{M=2}$  (2-Categories)

$c_0 \xrightarrow{c_1} c_0'$  s-invertible if

$$\exists c_1^*: c_0' \rightarrow c_0, \exists \eta: 1_{c_0} \rightarrow c_1 c_1^*, \varepsilon: c_1^* c_1 \rightarrow 1_{c_0'}$$

$$\text{and } \exists \eta^*: c_1^* c_1 \rightarrow 1_{c_0}, \varepsilon^*: 1_{c_0'} \rightarrow c_1^* c_1$$

$$\text{s.t. } \eta \eta^* = 1_{1_{c_0}}$$

$$\varepsilon \varepsilon^* = 1_{c_1^* c_1}$$

$$\eta^* \eta = 1_{c_1 c_1^*}$$

$$\varepsilon^* \varepsilon = 1_{1_{c_0'}}$$

→ equiv

II) [M.M. Kapranov, V.A. Voevodsky, 1991]

give a definition of  $\omega$ -groupoid as n-category with weak inverses.

SLOGAN

All equations

$$a \cdot x \cong b$$

admit a (weak) solution.

(--- and  $x \cdot a \cong b$  also)

THEN

a solution of  $a \cdot x = 1$

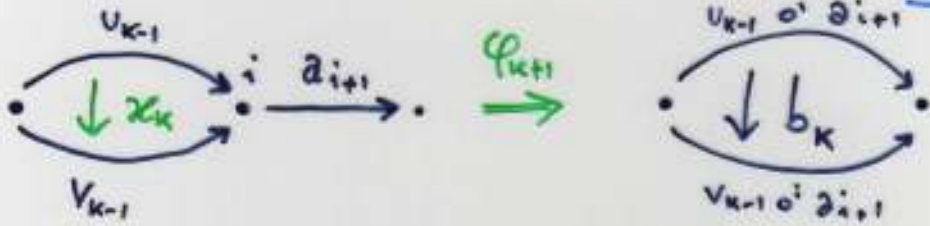
gives a possible inverse of  $a$ .

DEFINITION (KV)  $\mathcal{C}$  is a

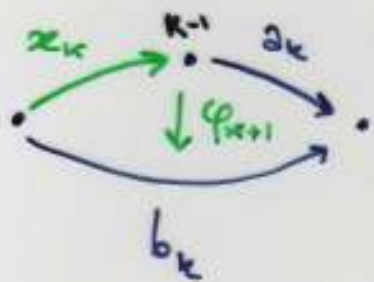
$n$ -groupoid if  $\forall i < k \leq n$

$GR'_{i,k} \forall a_{i+1}, b_k, u_{k-1}, v_{k-1}$ , (with  $i < k-1$ ) below

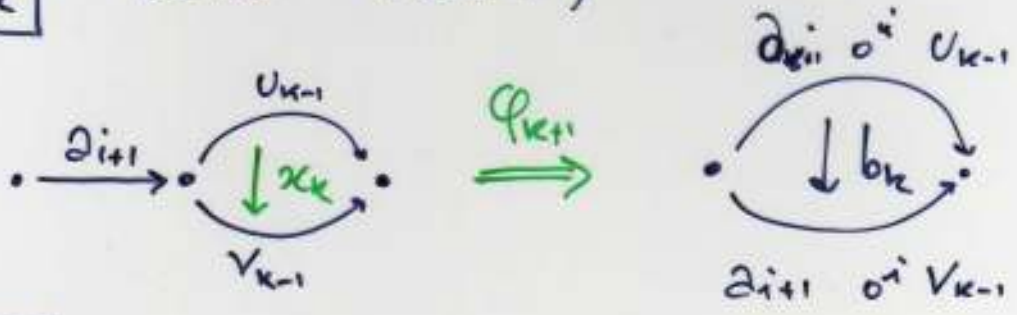
$\exists x_k : u_{k-1} \rightarrow v_{k-1}$ ,  $\varphi_{k+1} : x_k \circ^i a_{i+1} \rightarrow b_k$



$GR'_{k-1,k} \forall a_k, b_k \exists x_k, \varphi_{k+1}$



$GR''_{i,k}$  (with  $i < k-1$ )



$GR''_{k-1,k}$

A commutative diagram with four objects. The top-left object is  $\cdot$ , the top-right is  $\cdot$ , the bottom-left is  $\cdot$ , and the bottom-right is  $b_k$ . A curved arrow labeled  $a_k$  goes from  $\cdot$  to  $\cdot$ . A curved arrow labeled  $x_k$  goes from  $\cdot$  to  $b_k$ . A curved arrow labeled  $b_k$  goes from  $\cdot$  to  $b_k$ . A curved arrow labeled  $\varphi_{k+1}$  goes from  $\cdot$  to  $b_k$ . A vertical arrow labeled  $\downarrow \varphi_{k+1}$  connects  $\cdot$  to  $b_k$ .

[K.V.]  $n$ -groupoid  $\Rightarrow$  [S]  $n$ -groupoid

i.e.  $\boxed{GR'_{i,k}}, \boxed{GR'_{k-1,k}}, \boxed{GR''_{i,k}}, \boxed{GR''_{k-1,k}} \Rightarrow \boxed{GR'_{k-1,k}}, \boxed{GR''_{k-1,k}}$

[K.V.] claim  ~~$\Leftarrow$~~

Their argument is:

[KV]-DEF  $\Rightarrow$  "coherent" system of  $g$ -inv.

and  $\ll$  there is no way to construct such a system from Street's def<sup>n</sup>.  $\gg$



### III) C.S.Q.I. (description in low dim)

$$\boxed{M=1} \quad c_1: C_0 \longrightarrow C_0', \quad \exists c_1^*: C_0' \longrightarrow C_0 \quad \text{s.t.}$$

$$c_1 c_1^* = 1_{C_0}, \quad c_1^* c_1 = 1_{C_0'}$$

$$\boxed{M=2} \quad \underline{c_1}: C_0 \longrightarrow C_0', \quad \exists \underline{c_1^*}: C_0' \longrightarrow C_0$$

$$\underline{p_2}: C_1 c_1^* \longrightarrow 1_{C_0}, \quad \lambda_2: c_1^* c_1 \longrightarrow 1_{C_0'}$$

s.t.

$$\begin{array}{ccc} & p_2 c_1 & \\ \curvearrowright & & \curvearrowleft \\ c_1 c_1^* c_1 & \xrightarrow{\quad} & c_1 \\ & \uparrow \text{green} & \\ & c_1 & \end{array}$$

$$\begin{array}{ccc} & \lambda_2 c_1^* & \\ \curvearrowright & & \curvearrowleft \\ c_1^* c_1 c_1^* & \xrightarrow{\quad} & c_1^* \\ & \downarrow \text{green} & \\ & c_1^* & \end{array}$$

$$(+ p_2^* \lambda_2^*)$$

$$c_1 \lambda_2$$

$$c_1^* p_2$$

$$\uparrow p_3$$

$$\downarrow \lambda_3$$

...

PROP [KV]  $\Rightarrow$  C.S.Q.I.

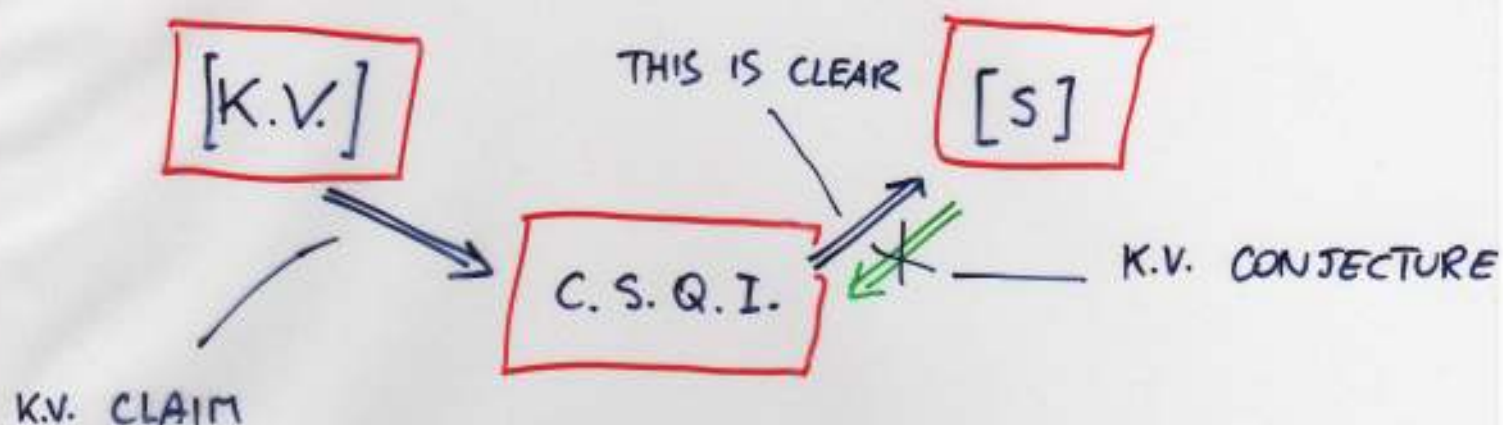
idea:

$$(c_1^*, p_2) \quad \text{solution of} \quad c_1 \circ \kappa \cong 1_{C_0}$$

$$(\lambda_2, p_3) \quad \text{solution of} \quad c_1 \circ \kappa \cong p_2 \circ c_1$$

....

The situation so far:



Need another notion of "inverse"

IV (variation on) [C. SIMPSON, 1998]

→ relies on a notion of equivalence of  $n$ -categories:

DEFINITION  $F: \mathcal{C} \rightarrow \mathcal{D}$   $n$ -equivalence

$n=0$  isomorphism of sets

$n > 0$  (E1)  $\forall c_0, c_0' \in \mathcal{C}_0$

$F_1^{c_0, c_0'}: \mathcal{C}_1(c_0, c_0') \rightarrow \mathcal{D}_1(Fc_0, Fc_0')$   $(n-1)$ -equiv.

(E2)  $\forall d_0 \in \mathcal{D}_0 \exists c_0 \in \mathcal{C}_0, d_1: d_0 \rightarrow Fc_0$

s.t.  $d_1 \circ \circ -$ ,  $- \circ \circ d_1$   $(n-1)$ -equiv.

small and...

Strict n-Categoriesinductive definition:

$$\mathcal{C} = \langle \mathcal{C}_0, \{ \mathcal{C}_1(a, a') \}_{a, a' \in \mathcal{C}_0} \rangle$$

Set of objects

 $(n-1)$ -categories+ compositions and identities

$$\forall a, a', a'' \in \mathcal{C}_0, \quad \circ: \mathcal{C}_1(a, a') \times \mathcal{C}_1(a', a'') \longrightarrow \mathcal{C}_1(a, a'')$$

$$\forall a \in \mathcal{C}_0, \quad U(a): \mathbb{1} \longrightarrow \mathcal{C}_1(a, a)$$

+ axioms: $(n-1)$ -functors satisfying

associativity, left/right unit

axioms.

Strict n-functors ...

DEFINITION

$$c_1 : c_0 \rightarrow c_0'$$

is w-invertible if  $\forall c_0'' \in \mathcal{C}$

$$(i) \quad c_1 \circ - : \mathcal{C}_1(c_0', c_0'') \longrightarrow \mathcal{C}_1(c_0, c_0'')$$

$$(ii) \quad - \circ c_1 : \mathcal{C}_1(c_0'', c_0) \longrightarrow \mathcal{C}_1(c_0'', c_0')$$

are equivalences of (n-1)-categories

(? (i)  $\Leftrightarrow$  (ii) ?)

DEFINITION

$\mathcal{C}$  is a n-groupoid:  $\boxed{n=0}$   $\mathcal{C}$  is a set

$\boxed{n>0}$

a)  $\forall c_0, c_0' \quad \mathcal{C}_1(c_0, c_0')$  is a (n-1)-groupoid

b)  $\forall c_1 : c_0 \rightarrow c_0'$  is w-invertible.

**THEOREM** (Simpson)

$\mathcal{C}$  (K.V.)-n-groupoid  $\Leftrightarrow$  (Simpson)-n-groupoid

in particular,

$c_1: C_0 \rightarrow C_0'$  [KV]-invertible  $\Leftrightarrow$   $\omega$ -invertible

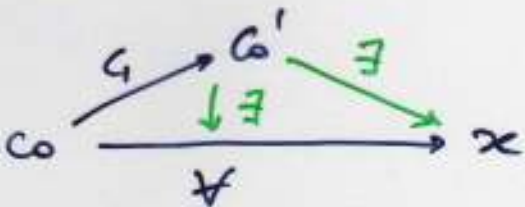
$n=1$   $c_1$   $\omega$ -invert.  $\stackrel{\text{def}}{\Leftrightarrow}$   $\begin{cases} C_0 - \\ - \circ c_1 \end{cases}$  isos of sets  
 $\Leftrightarrow c_1$  is an iso of  $\mathcal{C}$

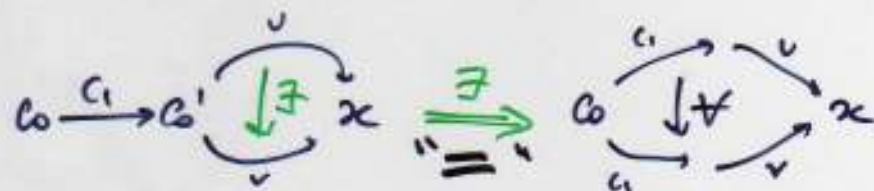
$n=2$   $c_1$   $\omega$ -invert.  $\stackrel{\text{def}}{\Leftrightarrow}$   $\begin{cases} C_0 - \\ - \circ c_1 \end{cases}$  equiv. of cats.

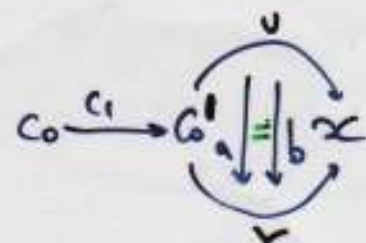
i.e.

(i)  $\forall x \in C_0 \quad C_0 - : \mathcal{C}(C_0', x) \rightarrow \mathcal{C}(C_0, x)$

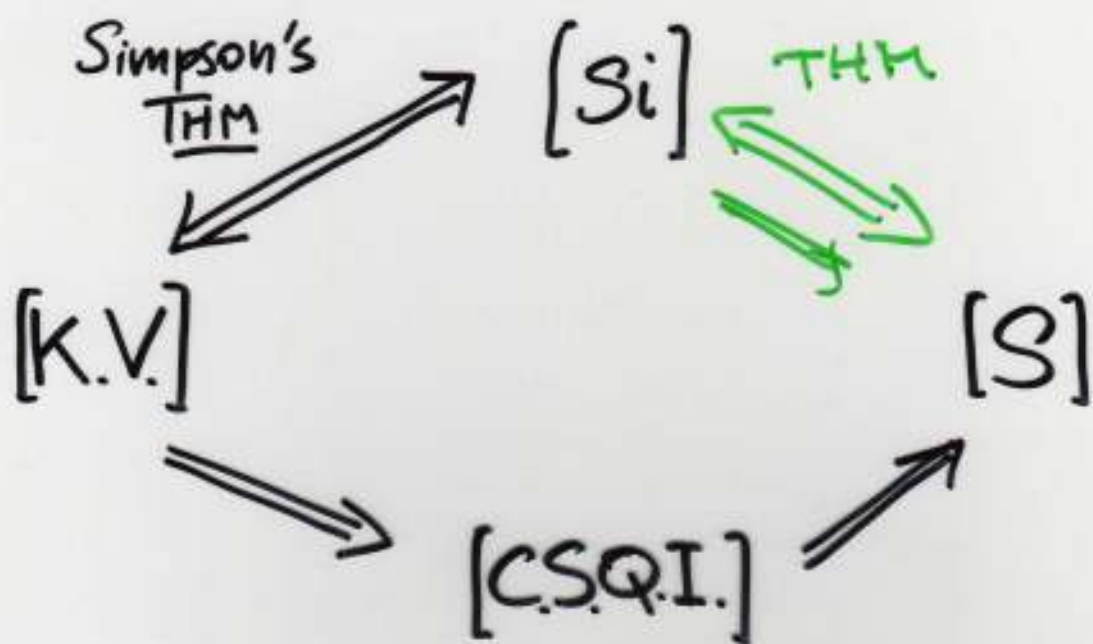
is an equiv. of cats when

•) ess. surjective : 

•) full 

•) faithful   $\forall " =_1 " : c_1 a \rightarrow c_1 b$   
 $\exists " =_2 " a \rightarrow b, " =_1 " \Rightarrow " =_2 "$

+ duals (ii)



## THEOREM (Kasangian, M, Vitale)

A 1-cell  $c_1 : c_0 \rightarrow c_0'$  of a  $n$ -cat  $\mathcal{C}$  is  $w$ -invertible iff  $s$ -invertible

## COROLLARY

The corresponding notions of  $n$ -groupoid coincide.

The proof is not straightforward  
 (... still looking for a better one!)

Key - idea :

equivalence of  $n$ -categories =  $n+1$  - levels of surjectivity

for instance,  $n=1$  (1-cat = cat)

a functor is an equivalence if

ess. surjective = "surjective" on objects

full = "surjective" on arrows

faithful = "surjectiv" on equivs

LOCALIZE  
 ↓

Main tool : inductive def

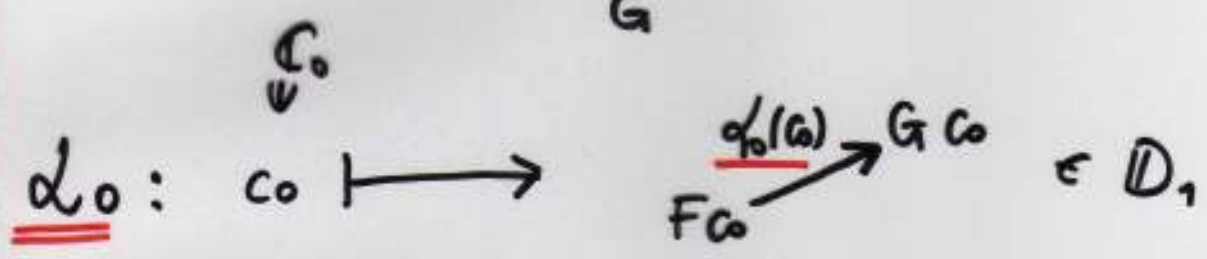
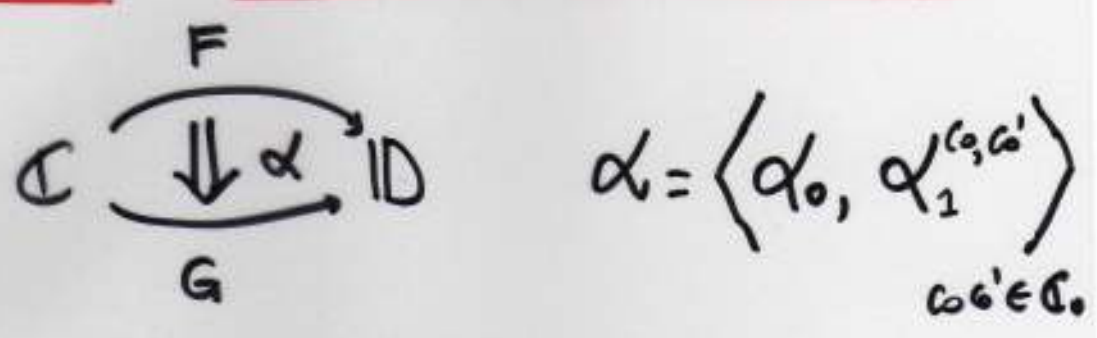
Consider n-Cat as a sesquicategory

obj : n-categories

mor : n-functors

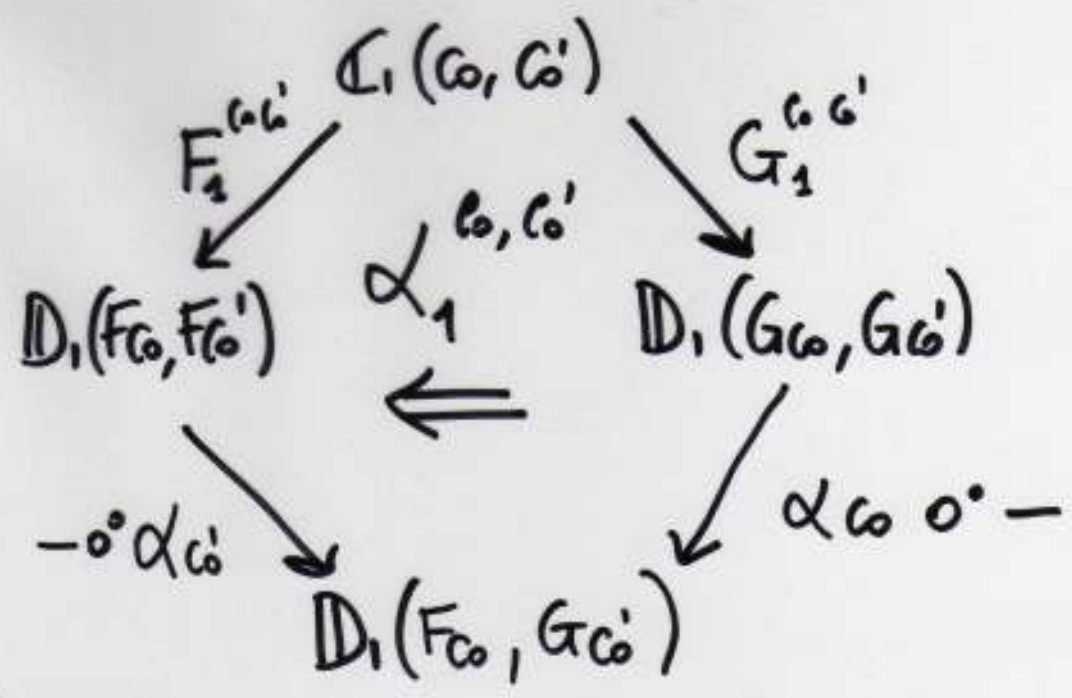
2-mor : weak n-transformations

Defn



$\forall c_0, c'_0 \in C_0,$

$\alpha_1^{(c_0, c')}$



+ axioms



(idea of the) proof:

We prove that  $C_1 \circ \dots \circ$

is an equivalence of  $(n-1)$ -categories

by double induction over

- the dimension

- the "level" of localization

# Case Study : direct

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## Proof for $n=2$

Let  $c_1$  be  $S$ -invertible, i.e.

$$\exists c_1^*, \eta: 1_{C_0} \rightarrow c_1 c_1^*, \eta^* \text{ s.t. } \eta \eta^* = 1, \eta^* \eta = 1$$
$$\varepsilon: c_1^* c_1 \rightarrow 1_{C_0'}, \varepsilon^* \text{ s.t. } \varepsilon^* \varepsilon = 1, \varepsilon \varepsilon^* = 1$$

we want to show  $c_1 \circ -$ ,  $- \circ c_1$  equiv.

(i) fix  $c_0''$  and consider

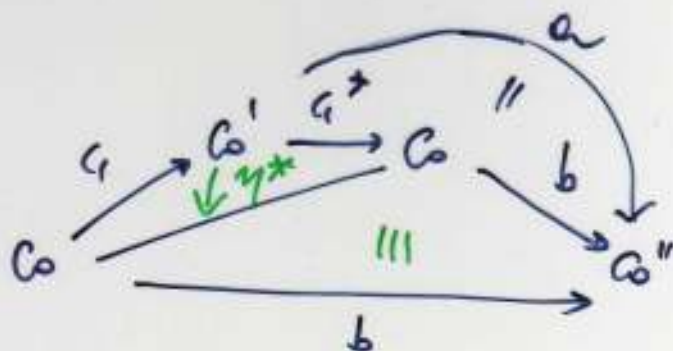
$$c_1 \circ -: \mathcal{C}(C_0', C_0'') \longrightarrow \mathcal{C}(C_0, C_0'')$$

• ess. surjective

$\forall b: C_0 \rightarrow C_0'', \exists a: C_0' \rightarrow C_0''$  s.t.  $c_1 a = b$ ?

let  $a = c_1^* b$

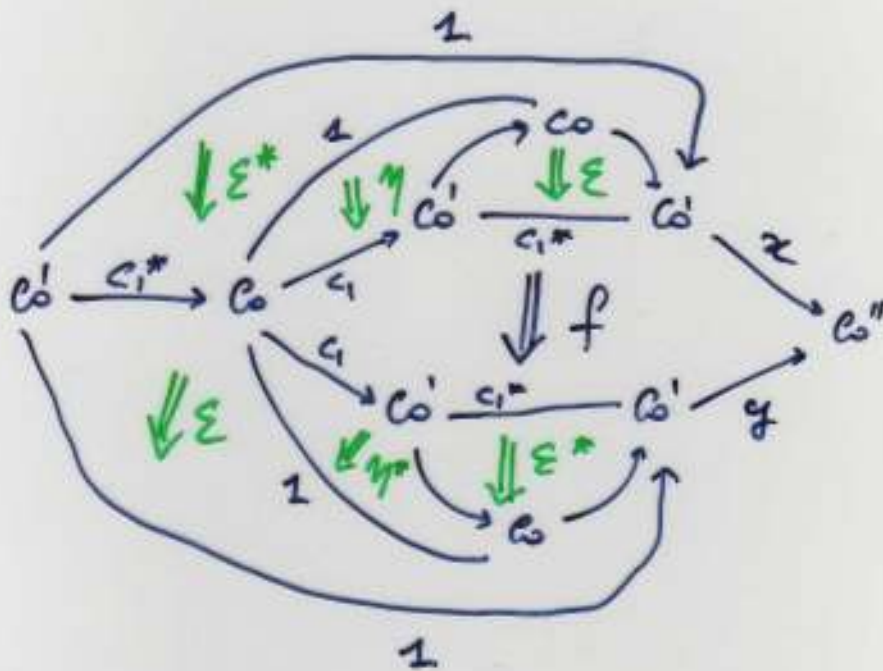
then



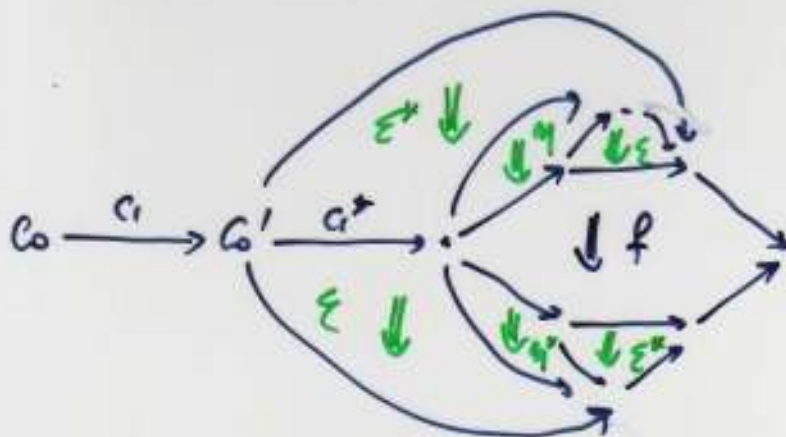
• full : let  $x, y : C_0' \rightarrow C_0''$ ; 16

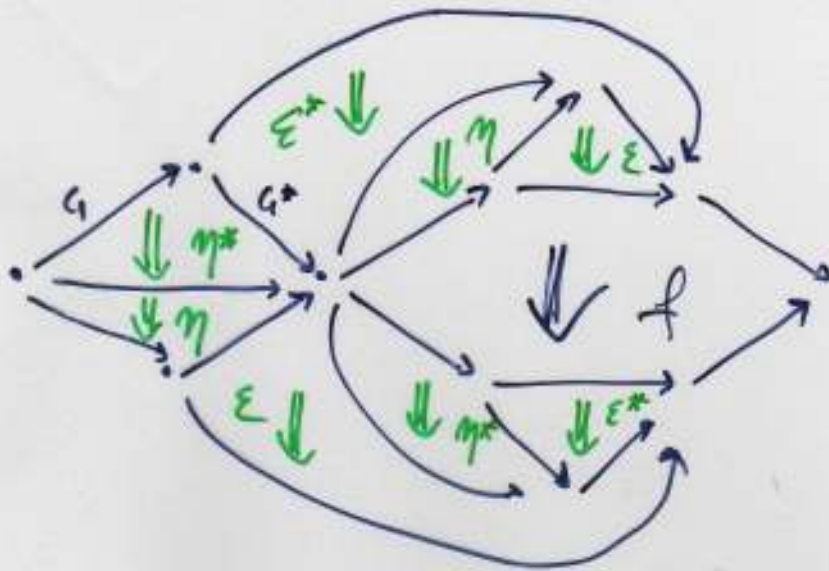
$$\forall C_0 \begin{array}{c} \xrightarrow{c_1 x} \\ \downarrow f \\ \xrightarrow{c_1 y} \end{array} C_0'' \quad \exists C_0' \begin{array}{c} \xrightarrow{x} \\ \downarrow g \\ \xrightarrow{y} \end{array} C_0'' \quad \text{s.t. } c_1 g = f ?$$

let  $g$  be the following pasting:

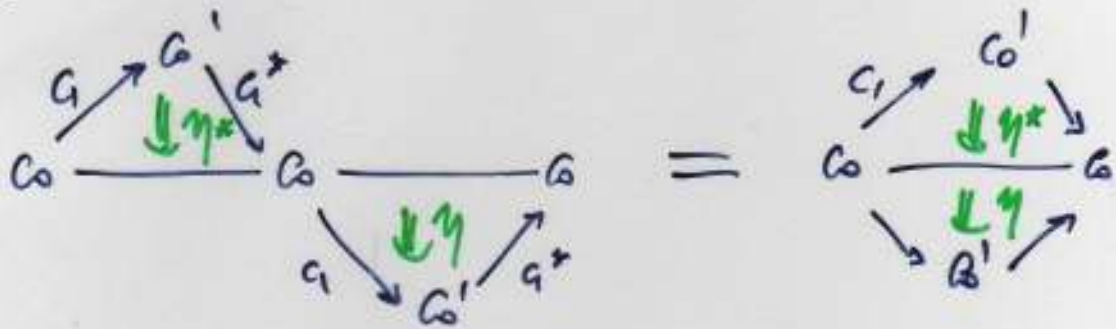


compute  $c_1 \circ g$  :

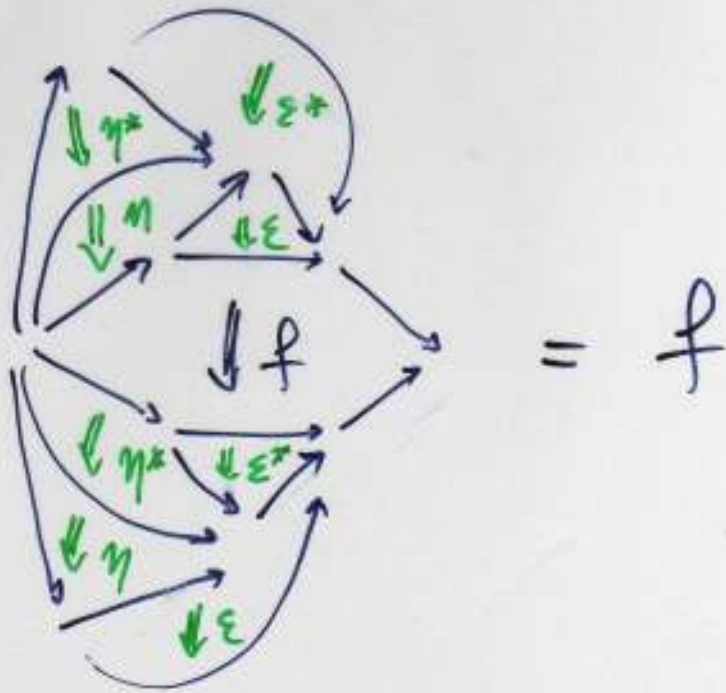




But



hence



• faithful: let  $x, y: \mathcal{C}' \rightarrow \mathcal{C}''$

$$\forall \begin{array}{ccc} \mathcal{C}' & \begin{array}{c} \xrightarrow{x} \\ \downarrow f \\ \downarrow g \\ \xrightarrow{y} \end{array} & \mathcal{C}'' \\ & \text{if } c_1 f = a g & \Rightarrow f = g? \end{array}$$

consider

$$f = \begin{array}{ccc} & \downarrow \varepsilon & \\ \mathcal{C}' & \xrightarrow{c_1^*} \mathcal{C} & \xrightarrow{c_1} \mathcal{C}' \\ & \downarrow \varepsilon^* & \\ & & \end{array} \begin{array}{ccc} \mathcal{C}' & \begin{array}{c} \xrightarrow{x} \\ \downarrow f \\ \downarrow g \\ \xrightarrow{y} \end{array} & \mathcal{C}'' \\ & \text{if } c_1 f = a g & \Rightarrow f = g? \end{array} =$$

$$= \begin{array}{ccc} \begin{array}{ccc} \xrightarrow{\quad} & \downarrow \varepsilon & \xrightarrow{\quad} \\ \xrightarrow{\quad} & \xrightarrow{\quad} & \xrightarrow{\quad} \end{array} & \begin{array}{ccc} \xrightarrow{\quad} & \downarrow f & \xrightarrow{\quad} \\ \xrightarrow{\quad} & \xrightarrow{\quad} & \xrightarrow{\quad} \end{array} & = \\ \begin{array}{ccc} \xrightarrow{\quad} & \downarrow \varepsilon^* & \xrightarrow{\quad} \\ \xrightarrow{\quad} & \xrightarrow{\quad} & \xrightarrow{\quad} \end{array} & & \end{array}$$

$$= \begin{array}{ccc} \begin{array}{ccc} \xrightarrow{\quad} & \downarrow \varepsilon & \xrightarrow{\quad} \\ \xrightarrow{\quad} & \xrightarrow{\quad} & \xrightarrow{\quad} \end{array} & \begin{array}{ccc} \xrightarrow{\quad} & \downarrow g & \xrightarrow{\quad} \\ \xrightarrow{\quad} & \xrightarrow{\quad} & \xrightarrow{\quad} \end{array} & = g \\ \begin{array}{ccc} \xrightarrow{\quad} & \downarrow \varepsilon^* & \xrightarrow{\quad} \\ \xrightarrow{\quad} & \xrightarrow{\quad} & \xrightarrow{\quad} \end{array} & & \end{array}$$