

When is $\mathbf{Cat}(\mathcal{Q})$ Cartesian closed?

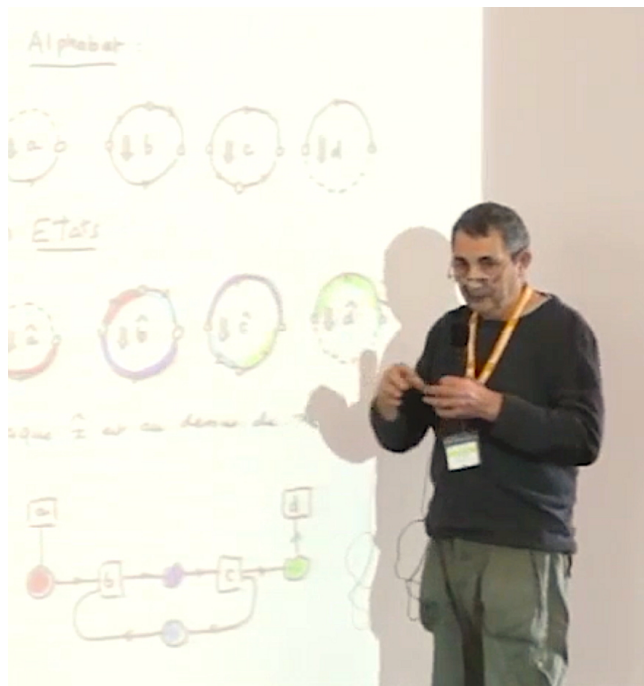
Isar Stubbe

joint work with Junche Yu

Université du Littoral, France

Lens, 6 June 2025

Mais d'abord...

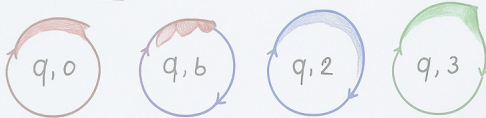


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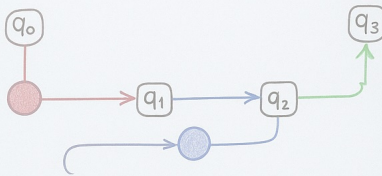
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Etats



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Exponentials, Cartesian closedness

Let \mathcal{C} be a category with finite products.

An object $A \in \mathcal{C}$ is **exponentiable** when

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This means: for any $B \in \mathcal{C}$ there exists an object B^A and a morphism $e_B: B^A \times A \rightarrow B$ in \mathcal{C} inducing, for any $X \in \mathcal{C}$, a (natural) bijection

$$\mathcal{C}(X \times A, B) \cong \mathcal{C}(X, B^A)$$

The diagram is a commutative triangle with vertices $X \times A$, $B^A \times A$, and B . A horizontal arrow points from $X \times A$ to B and is labeled $\forall g$. A diagonal arrow points from $X \times A$ down to $B^A \times A$ and is labeled $\exists! g' \times A$. A diagonal arrow points from $B^A \times A$ up to B and is labeled e_B .

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A commutative triangle diagram with vertices $X \times A$ (top-left), $B^A \times A$ (bottom), and B (top-right).
- A horizontal arrow from $X \times A$ to B is labeled $\forall g$.
- A diagonal arrow from $X \times A$ to $B^A \times A$ is labeled $\exists! g' \times A$.
- A diagonal arrow from $B^A \times A$ to B is labeled e_B .

Putting $X = 1$ in the above, one has

$$\mathcal{C}(A, B) \cong \mathcal{C}(1, B^A),$$

so the “points” of the object B^A are precisely the morphisms from A to B ; and similarly $e_B: B^A \times A \rightarrow B$ is the “pointwise” evaluation of such morphisms.

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The category \mathcal{C} is **Cartesian closed** when every $A \in \mathcal{C}$ is exponentiable.

Example

Consider (A, \leq) and (B, \leq) in the category Ord of ordered sets and order-preserving maps.

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Because the “points” functor $\text{Ord}(1, -): \text{Ord} \rightarrow \text{Set}$ produces the underlying sets and maps of objects and morphisms in Ord , the underlying set and evaluation map of their powerobject must be

$$B^A = \text{Ord}(A, B) \quad \text{and} \quad e_B: B^A \times A \rightarrow B: (f, a) \mapsto fa.$$

But is there an order on B^A that makes e_B order-preserving and universal?

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To make the evaluation map order-preserving we must have

$$\text{if } f \leq g \text{ in } B^A \text{ and } a \leq b \text{ in } A \text{ then } fa \leq gb \text{ in } B \quad \text{for all } f, g \in B^A \text{ and } a, b \in A,$$

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and its required universal property forces us to define

$$\begin{aligned} f \leq g \text{ in } B^A &\stackrel{\text{def}}{\iff} \text{for all } a, b \in A: \text{if } a \leq b \text{ then } fa \leq gb \\ &\iff \text{for all } a \in A: fa \leq ga \end{aligned}$$

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The category Ord is Cartesian closed.

Counterexample

Consider (A, d) and (B, d) in the category Met of generalized metric spaces and non-expansive maps.

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The evaluation is a non-expansive map iff

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Theorem (Clementino and Hofmann, 2006). *(A, d) is exponentiable in Met if and only if, for every $a, b \in A$, whenever $d(a, b) = r + s$ in $[0, \infty]$ then, for every $\varepsilon > 0$, there exists $m \in A$ such that $d(a, m) < r + \varepsilon$ and $d(m, b) < s + \varepsilon$.*

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The category Met is **not Cartesian closed**.

Whence our question



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In fact, both Ord and Met are of the form " $\text{Cat}(\mathcal{Q})$ ", the category of quantaloid-enriched categories and functors. One is Cartesian closed, the other is not. So Cartesian closedness of $\text{Cat}(\mathcal{Q})$ is a significant property of the quantaloid \mathcal{Q} .

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So exactly when is $\text{Cat}(\mathcal{Q})$ Cartesian closed?

Quantale-enriched categories

A **quantale** $Q = (Q, \bigvee, \circ, 1)$ is a completely ordered monoid in which

$$a \circ \left(\bigvee_i b_i \right) = \bigvee_i (a \circ b_i) \quad \text{and} \quad \left(\bigvee_i a_i \right) \circ b = \bigvee_i (a_i \circ b)$$

for all elements $a, b, (a_i)_i, (b_i)_i$ in Q . (So Q is a *very particular* \mathcal{V} .)

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A Q -**category** \mathbb{A} is

a set \mathbb{A}_0 of “objects”,

a function $\mathbb{A}: \mathbb{A}_0 \times \mathbb{A}_0 \rightarrow Q: (x, y) \mapsto \mathbb{A}(x, y)$ assigning “homs”,

such that, for any $x, y, z \in \mathbb{A}_0$,

$$\mathbb{A}(x, y) \circ \mathbb{A}(y, z) \leq \mathbb{A}(x, z) \quad \text{and} \quad 1 \leq \mathbb{A}(x, x).$$

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For $Q = (\{0, 1\}, \vee, \wedge, 1) : \mathbb{A}(x, y) = \llbracket x \leq y \rrbracket$ and $\text{Cat}(Q) = \text{Ord}$.

For $Q = ([0, \infty], \bigwedge, +, 0) : \mathbb{A}(x, y) = d(x, y)$ and $\text{Cat}(Q) = \text{Met}$.

But also t -norms and “fuzzy orders”, probabilistic metric spaces, monoidal topology, ...

Quantaloid-enriched categories

A **quantaloid** \mathcal{Q} is a (small) locally completely ordered bicategory in which

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a “typed” set $t: \mathbb{A}_0 \rightarrow \mathcal{Q}_0$ of “objects”,

a function $\mathbb{A}: \mathbb{A}_0 \times \mathbb{A}_0 \rightarrow \mathcal{Q}_1: (x, y) \mapsto \mathbb{A}(x, y)$ assigning “homs”,

such that, for any $x, y, z \in \mathbb{A}_0$,

$$\mathbb{A}(x, y) \in \mathcal{Q}(ty, tx) \quad \text{and} \quad \mathbb{A}(x, y) \circ \mathbb{A}(y, z) \leq \mathbb{A}(x, z) \quad \text{and} \quad 1 \leq \mathbb{A}(x, x).$$

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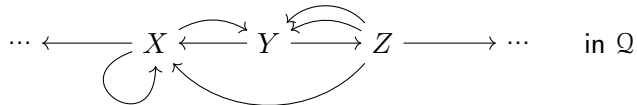
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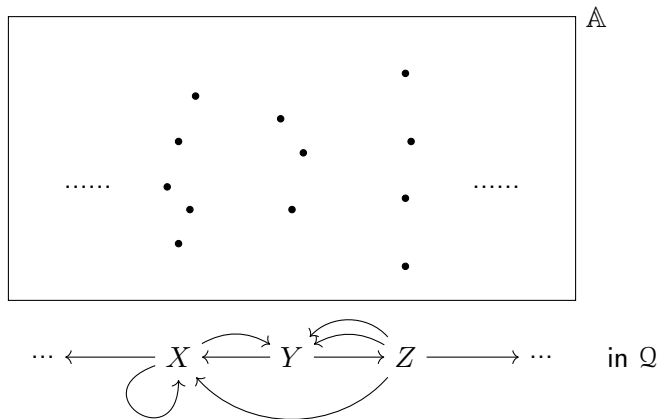
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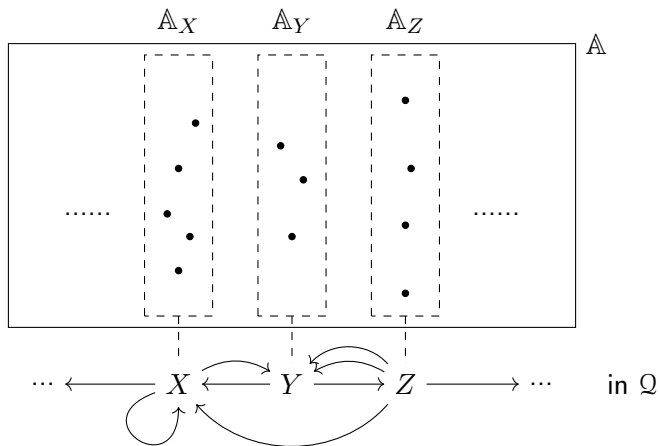
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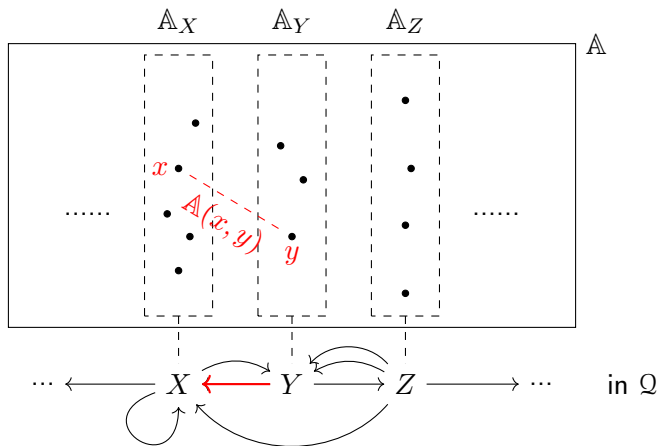
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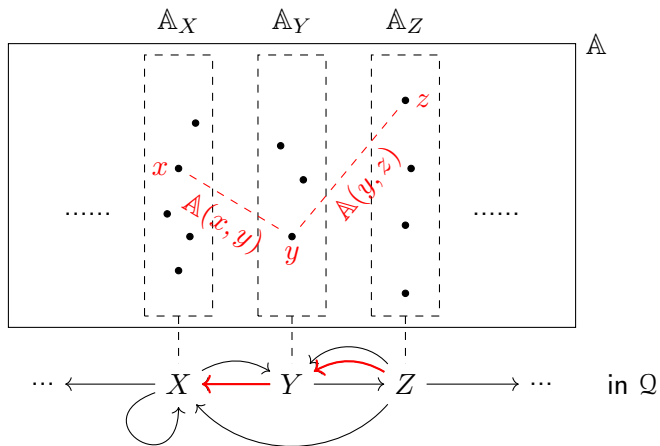
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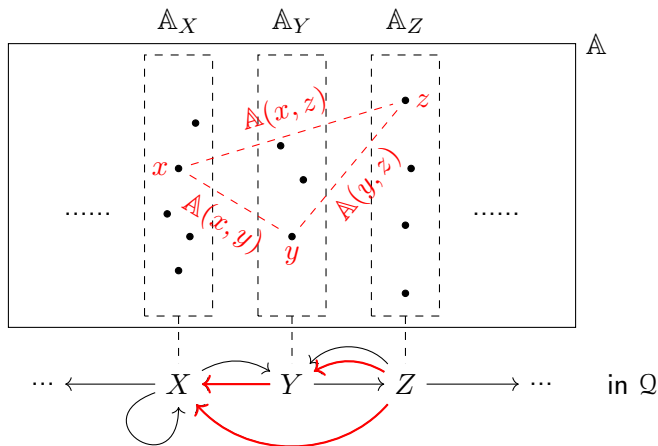
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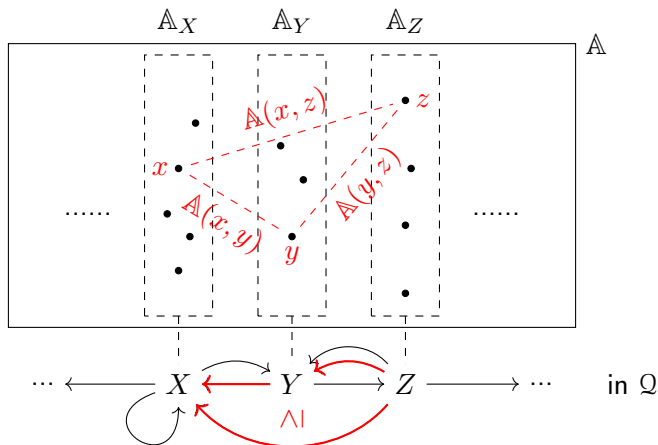
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A **quantale** is exactly a one-object quantaloid (think of “group” vs. “groupoid”), so we recover all previous examples.

Performing universal constructions on a quantale very often produce a quantaloid. These are used in many new examples: partial (probabilistic) metric spaces, sheaves, ...

Exponentiability in $\text{Cat}(\mathcal{Q})$

The product $\mathbb{A} \times \mathbb{B}$ of two \mathcal{Q} -categories \mathbb{A} and \mathbb{B} exists in $\text{Cat}(\mathcal{Q})$:

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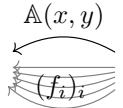
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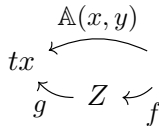
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Theorem (Clementino, Hofmann and Stubbe, 2009). *A \mathcal{Q} -category \mathbb{A} is exponentiable in $\text{Cat}(\mathcal{Q})$ if and only if*

$$1. \text{ for all } tx \overset{\mathbb{A}(x, y)}{\curvearrowright} ty : \left(\bigvee_i f_i \right) \wedge \mathbb{A}(x, y) = \bigvee_i \left(f_i \wedge \mathbb{A}(x, y) \right),$$


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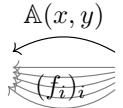
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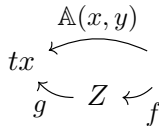
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The first condition says that "every hom of \mathbb{A} is exponentiable"; the second condition says that " \mathbb{A} has enough intermediate objects".

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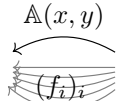
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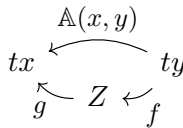
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Necessity of the conditions is checked by "probing" exponentiability of \mathbb{A} on particular \mathbb{B} 's. \square

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By definition, $\text{Cat}(\mathcal{Q})$ is Cartesian closed when **all** \mathcal{Q} -categories are exponentiable.

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homs: $\mathbb{A}_h(*_1, *_2) = h$, $\mathbb{A}_h(*_2, *_1) = 0_{Y,X}$, $\mathbb{A}_h(*_1, *_1) = 1_Y$ and $\mathbb{A}_h(*_2, *_2) = 1_X$

Cartesian closedness of $\text{Cat}(\mathcal{Q})$

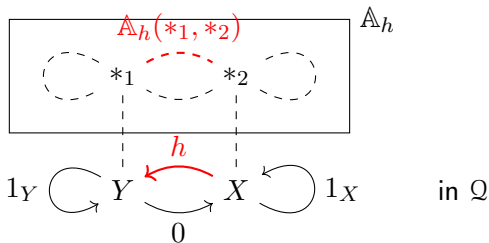
By definition, $\text{Cat}(\mathcal{Q})$ is Cartesian closed when **all** \mathcal{Q} -categories are exponentiable.

So we can find **necessary conditions** for Cartesian closedness of $\text{Cat}(\mathcal{Q})$ by requiring exponentiability of **some** \mathcal{Q} -categories.

For any $h: X \rightarrow Y$ in \mathcal{Q} , let \mathbb{A}_h be the \mathcal{Q} -category defined as:

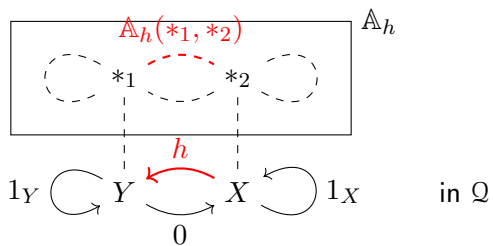
objects: $(\mathbb{A}_h)_0 = \{*_1, *_2\}$ with types $t*_1 = Y$ and $t*_2 = X$,

homs: $\mathbb{A}_h(*_1, *_2) = h$, $\mathbb{A}_h(*_2, *_1) = 0_{Y,X}$, $\mathbb{A}_h(*_1, *_1) = 1_Y$ and $\mathbb{A}_h(*_2, *_2) = 1_X$

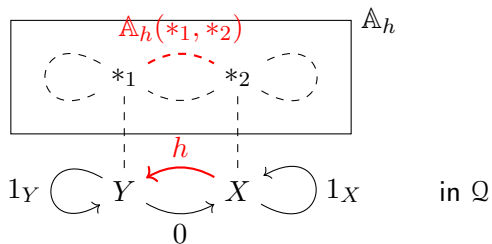


It is the *collage* (= universal cotabulation) of the image of h under the inclusion $\mathcal{Q} \rightarrow \text{Dist}(\mathcal{Q})$.

Cartesian closedness of $\text{Cat}(\mathcal{Q})$



Cartesian closedness of $\text{Cat}(\mathcal{Q})$

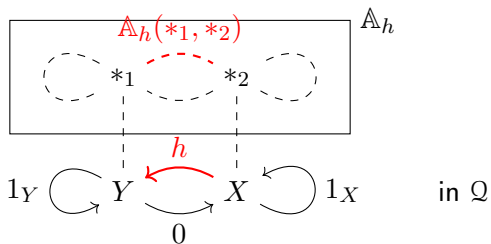


Theorem. \mathbb{A}_h is exponentiable if and only if

1. for all $tx \xrightarrow{\mathbb{A}_h(x,y)} ty$: $\left(\bigvee_i f_i \right) \wedge \mathbb{A}_h(x,y) = \bigvee_i \left(f_i \wedge \mathbb{A}_h(x,y) \right)$,

2. for all $tx \xrightarrow{\mathbb{A}_h(x,y)} ty$: $(g \circ f) \wedge \mathbb{A}_h(x,y) = \bigvee_{z \in (\mathbb{A}_h)_Z} \left(g \wedge \mathbb{A}_h(x,z) \right) \circ \left(f \wedge \mathbb{A}_h(z,y) \right)$.

Cartesian closedness of $\text{Cat}(\mathcal{Q})$

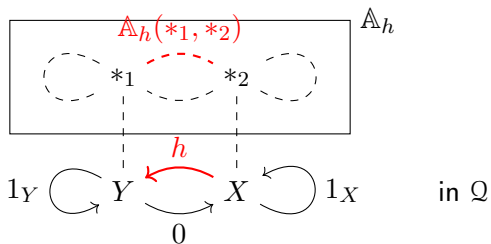


Theorem. \mathbb{A}_h is exponentiable ~~if and~~ only if, for $x = *_1$ and $y = *_2$,

$$1. \text{ for all } tx \overset{\mathbb{A}_h(x, y)}{\curvearrowright} ty : \left(\bigvee_i f_i \right) \wedge \mathbb{A}_h(x, y) = \bigvee_i \left(f_i \wedge \mathbb{A}_h(x, y) \right),$$

$$2. \text{ for all } tx \overset{\mathbb{A}_h(x, y)}{\curvearrowright} ty : (g \circ f) \wedge \mathbb{A}_h(x, y) = \bigvee_{z \in (\mathbb{A}_h)_Z} \left(g \wedge \mathbb{A}_h(x, z) \right) \circ \left(f \wedge \mathbb{A}_h(z, y) \right).$$

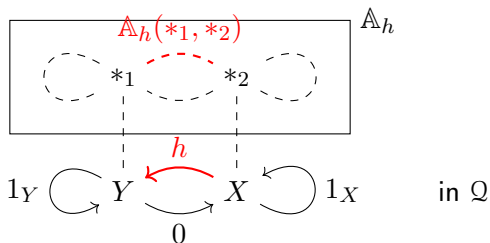
Cartesian closedness of $\text{Cat}(\mathcal{Q})$



Theorem. \mathbb{A}_h is exponentiable ~~if and~~ only if

- for all $Y \begin{array}{c} \xleftarrow{h} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \xleftarrow{(f_i)_i} \end{array} X : \left(\bigvee_i f_i \right) \wedge h = \bigvee_i (f_i \wedge h),$
- for all $Y \begin{array}{c} \xleftarrow{h} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \\ \xleftarrow{g} \quad Z \quad \xleftarrow{f} \end{array} X : (g \circ f) \wedge h = \bigvee_{z \in (\mathbb{A}_h)_Z} \left(g \wedge \mathbb{A}_h(*_1, z) \right) \circ \left(f \wedge \mathbb{A}_h(z, *_2) \right).$

Cartesian closedness of $\text{Cat}(\mathcal{Q})$

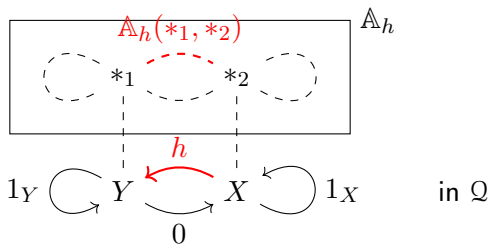


Theorem. \mathbb{A}_h is exponentiable ~~if and~~ only if

1. $- \wedge h: \mathcal{Q}(X, Y) \rightarrow \mathcal{Q}(X, Y)$ preserves suprema,

2. for all $Y \xleftarrow{h} X$: $(g \circ f) \wedge h = \bigvee_{z \in (\mathbb{A}_h)Z} (g \wedge \mathbb{A}_h(*_1, z)) \circ (f \wedge \mathbb{A}_h(z, *_2))$.

Cartesian closedness of $\text{Cat}(\mathcal{Q})$



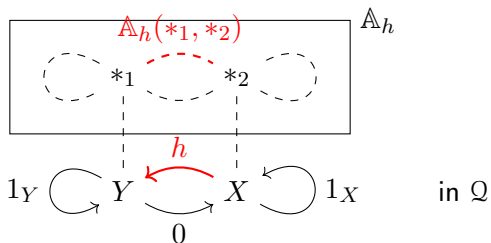
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$$= \begin{cases} 0_{X,Y} & \text{if } X \neq Z \neq Y \text{ (so } (\mathbb{A}_h)_Z = \emptyset), \\ (g \wedge h) \circ (f \wedge 1_X) & \text{if } X = Z \neq Y \text{ (so } (\mathbb{A}_h)_Z = \{*_2\}), \\ (g \wedge 1_Y) \circ (f \wedge h) & \text{if } X \neq Z = Y \text{ (so } (\mathbb{A}_h)_Z = \{*_1\}), \\ ((g \wedge h) \circ (f \wedge 1_X)) \vee ((g \wedge 1_Y) \circ (f \wedge h)) & \text{if } X = Z = Y \text{ (so } (\mathbb{A}_h)_Z = \{*_1, *_2\}). \end{cases}$$

Cartesian closedness of $\text{Cat}(\mathcal{Q})$



Theorem. Every \mathbb{A}_h is exponentiable ~~if and~~ only if

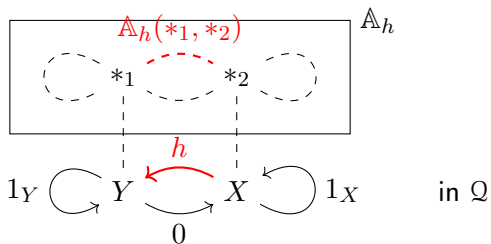
1. every $\mathcal{Q}(X, Y)$ is localic,

2. for all $X \xleftarrow{h} Y$:

$$\begin{array}{c} X \xleftarrow{h} Y \\ \quad \swarrow \quad \nwarrow \\ \quad Z \xleftarrow{f} \\ \quad \swarrow \quad \nwarrow \\ \quad X \quad Y \end{array}$$

$$(g \circ f) \wedge h = \begin{cases} 0_{X,Y} & \text{if } X \neq Z \neq Y, \\ (g \wedge h) \circ (f \wedge 1_X) & \text{if } X = Z \neq Y, \\ (g \wedge 1_Y) \circ (f \wedge h) & \text{if } X \neq Z = Y, \\ \left((g \wedge h) \circ (f \wedge 1_X) \right) \vee \left((g \wedge 1_Y) \circ (f \wedge h) \right) & \text{if } X = Z = Y. \end{cases}$$

Cartesian closedness of $\text{Cat}(\mathcal{Q})$



Theorem. $\text{Cat}(\mathcal{Q})$ is Cartesian closed if and only if

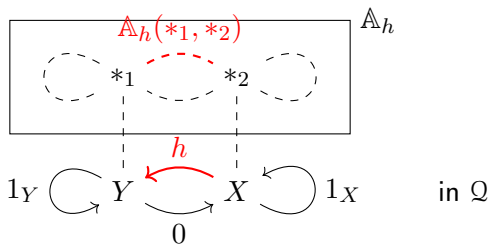
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Proof: “case analysis” for the exponentiability of any \mathcal{Q} -category. □

More examples

For a quantale $Q = (Q, \circ, 1)$, $\text{Cat}(Q)$ is Cartesian closed if and only if the underlying suplattice of Q is a locale and

$$\text{for all } a, b, c \in Q: (a \circ b) \wedge c = \left((a \wedge c) \circ (b \wedge 1) \right) \vee \left((1 \wedge a) \circ (b \wedge c) \right),$$

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If Q is an integral quantale (meaning that $1 = \top$), this further simplifies to

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Some quantale examples:

1. whenever $a \circ b = a \wedge b$ in Q (i.e. \mathcal{Q} is a locale),
2. for $Q = \{0, \frac{1}{2}, 1\}$ with natural order and multiplication $x \circ y = \max\{x + y - 1, 0\}$ (in [Lai and Zhang, 2016] with an *ad hoc* proof),
3. for $\mathcal{Q} = [0, 1]$ with natural order and multiplication $x \circ y = \begin{cases} 0 & \text{if } x, y \leq \frac{1}{2} \\ x \wedge y & \text{otherwise} \end{cases}$ (this is a *left-continuous t-norm*),
4. the only *continuous t-norm* satisfying the above condition is the Gödel *t-norm*, i.e. $\mathcal{Q} = [0, 1]$ with natural order and $x \circ y = x \wedge y$ (see [Lai and Zhang, 2016]).

More examples

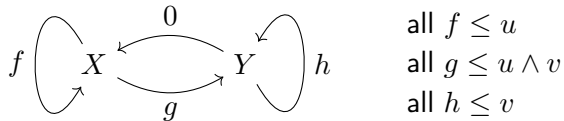
Some quantaloidal examples:

1. The conditions in the [Theorem](#) are stable under coproducts (but not under splitting of idempotents nor the construction of diagonals), so any coproduct of quantales satisfying the conditions is a quantaloid satisfying the conditions

More examples

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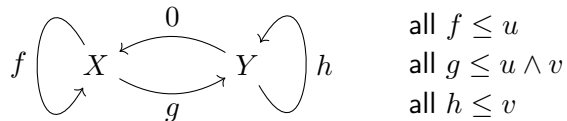
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3. The free quantaloid \mathcal{PC} on a (small) category \mathcal{C} is given by:

$$(\mathcal{PC})_0 = \mathcal{C}_0,$$

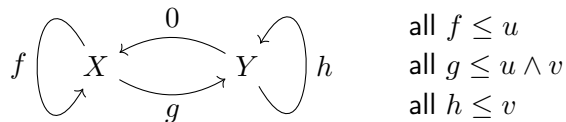
$$\mathcal{PC}(X, Y) = \mathcal{P}(\mathcal{C}(X, Y)) \text{ with } \bigcup \text{ as suprema,}$$

$$1_X = \{1_X\} \text{ and } G \circ F = \{g \circ f \mid g \in G, f \in F\}.$$

More examples

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3. The free quantaloid $\mathcal{P}\mathcal{C}$ on a (small) category \mathcal{C} is given by:

$$(\mathcal{P}\mathcal{C})_0 = \mathcal{C}_0,$$

$$\mathcal{P}\mathcal{C}(X, Y) = \mathcal{P}(\mathcal{C}(X, Y)) \text{ with } \bigcup \text{ as suprema,}$$

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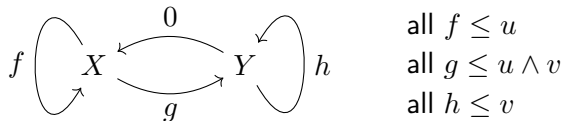
$\mathcal{P}\mathcal{C}$ is always locally localic, and satisfies the conditions in the [Theorem](#) if and only if

when two morphisms compose in \mathcal{C} , then at least one of them is an identity.

More examples

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when two morphisms compose in \mathcal{C} , then at least one of them is an identity.

In particular, the only free quantale $\mathcal{P}M$ (on a monoid $(M, \circ, 1)$) satisfying this condition is when $M = \{*\}$ (and so $\mathcal{P}M = (\{0, 1\}, \vee, \wedge, 1)$).

(This corrects a mistake in an example in [\[Clementino, Hofmann and Stubbe, 2009\]](#).)

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